



Fast Algorithm for Modularity-based Graph Clustering

Hiroaki Shiokawa NTT Software Innovation Center, NTT Corporation, July 23rd, 2013

© 2013 NTT Software Innovation Center

BACKGROUND & MOTIVATION

Large Graphs

- Large-scale graphs become available
 - Facebook: 1.11 billion active users / month(*1)
 - Twitter: <u>140 million active users</u> / day
 <u>340 million new posts</u> / day (*2)
 - And more …

• A lot of techniques for analyzing massive-scale graph

- Massive data require so much time for analysis ☺
- \bullet It is important to analyze large scale data quickly $\ensuremath{\textcircled{\scale}}$

(*1) "Key Facts", http://newsroom.fb.com/Key-Facts (*2) http://dev.twitter.com/media/authors

© 2013 NTT Software Innovation Center

Graph Clustering

• Graph clustering is one of the most important methods

- Community detection over social networks
- Event detection from microblogging services
- Brain Analysis, Image segmentation, ...



Modularity-based Graph Clustering

- Clustering methods which find the division of graph to maximize the modularity measure
- Improvement of clustering speed



Objective and Contributions

Objective

Fast graph clustering method with high modularity

3 key techniques

- 1. Incremental nodes aggregation
- 2. Incremental nodes pruning
- 3. Efficient ordering of nodes selections

Contributions of our algorithm

- Efficiency
 - Considerably faster than BGLL
 - Clusters 100M nodes within 3 minutes
- High Modularity
 - Scores high modularity as same as BGLL
- Effectiveness
 - Improves performances for complex networks

PRELIMINARIES

© 2013 NTT Software Innovation Center

Modularity

- Modularity evaluates the strength of division of a graph into clusters [Newman and Girvan 2004]
 - Finding the division which maximizes modularity is NP-complete
 - \Rightarrow A lot of greedy approaches were proposed



State of the art algorithm: BGLL

• Continuing following passes until the modularity score is maximized





Continue to the next pass

© 2013 NTT Software Innovation Center

PROPOSED ALGORITHM

Overview of proposed algorithm

Our method

- 3 key techniques
 - Incremental aggregation
 - Incremental pruning
 - Efficient ordering

BGLL

- Batch based aggregation
- Random ordering

- Clustering coefficient
- Power-law degree distribution





Idea 1 : Clustering coefficient

• Complex networks have large clustering coefficient

- Clustering coefficient is a measure of degree to which nodes in a graph tend to cluster together
- There are many duplicated nodes/edges in a graph which has large clustering coefficient



Idea 2 : Power-law degree distribution

• Complex networks have highly skewed degree distribution following the power-law distribution

- Most of nodes only have a few neighbor nodes, and only few nodes have large number of neighbor nodes
- The frequency F of nodes with degree d is $F \propto d^{-\alpha}$



 Random ordering of node selection leads redundant computation

Incremental Aggregation

- Incrementally aggregate nodes which belong to the same cluster
 - It aggregates duplicated edges between clusters while keeping the graph semantics
 - Example)



Incremental pruning

- Incrementally prune nodes whose cluster is trivially obtained
 - We can easily assign nodes to clusters without computing modularity gains
 - From the graph structure, there are 2 patterns of pruning



Incremental pruning (Cont.)

• Efficient pruning approach for pattern 2

- All nodes within the same cluster have been aggregated to a node by incremental aggregation
- We can find all prunable nodes by obtaining nodes such that they have only a single adjacent node



Efficient ordering of node selection

• Dynamically selects a node with the smallest degree

• Example) Node A and B being assigned to the same cluster



• By selecting node with the smallest degree, we can avoid producing super-cluster structures

EVALUATION

Datasets & Experimental Environment

Real world datasets

• 2 Social networks and 3 Web graphs of IP domains

Dataset		<i>E</i>	Skewness of degree distribution α
dblp-2010	326,186	1,615,400	2.82
ljournal-2008	5,363,260	79,023,142	2.29
uk-2005	39,459,925	936,364,282	1.71
webbase	118,142,155	1,019,903,190	2.14
uk2007-05	105,896,555	3,738,733,648	1.51

Experimental Environment

- All experiments were conducted on a Linux 2.6.18 server with Intel Xeon CPU L5640 2.27GHz and 144GB RAM
- Run all methods on 1 core, 1CPU

Computational time



Computational time – power-law differences



Computational time – size differences



Modularity score

Modularity score for datasets

- Large modularity score means the output of algorithms is well clustered
- Proposed method achieves <u>almost same modularity scores</u> <u>as/slightly higher than BGLL</u>

Table 2: Modularity Q

	db1p-2010	ljournal-2008	uk-2005	webbase-2001	uk-2007-05
Proposed	0.90	0.74	0.98	0.98	0.97
BGLL	0.88	0.74	0.97	0.96	0.97

CONCLUSION

Conclusion

• Fast clustering algorithm for large graphs

Our solution

- Incremental aggregation
- Incremental pruning
- Efficient ordering of nodes selections

Contribution of our algorithm

- Efficiency
 - Considerably faster than BGLL
 - Clusters 100M nodes within 3 minutes
- High Modularity
 - Scores high modularity as same as BGLL
- Effectiveness
 - Improves performance for complex networks